

Checkable Codes form Group Algebras to Group Rings

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Table of Contents

History

Group-Ring Codes

Code-Checkable Group Rings

References

History

- ▶ (MacWilliams, 1969)
"Codes and ideals in group algebras".
- ▶ (Hurley, 2007)
"Module codes over group rings".
- ▶ (Hurley, 2009)
"Codes from zero-divisors and units in group rings".
- ▶ (Jitman, 2010)
"Checkable Codes from group rings".

Codes Over Finite Fields

Let \mathbb{F}_q^n denote the vector space of all n -tuples over the finite field \mathbb{F}_q .

- ▶ An **(n,M) code** C over \mathbb{F}_q is a subset of \mathbb{F}_q^n of size M .
- ▶ If C is a k -dimensional vector subspace of \mathbb{F}_q^n , then C will be called **[n,k] linear code** over \mathbb{F}_q . This linear code C has q^k codewords.

Generator and Parity Check Matrix of a Linear Code

- ▶ A **generator matrix** for an $[n, k]$ linear code C is any $k \times n$ matrix G whose rows form a basis for the code C . C is written as:

$$C = \{xG : x \in \mathbb{F}_q^k\}$$

- ▶ A **parity check matrix** H for an $[n, k]$ linear code C is an $(n - k) \times n$ matrix defined by:

$$x \in C \Leftrightarrow Hx^T = 0$$

- ▶ Note that $GH^T = 0$. Thus G, H are in a sense zero-divisors.

Cyclic Codes

The class of cyclic codes is one of the most important classes of codes. In fact almost all codes used for practical issues, like BCH and Reed-Solomon codes, are cyclic codes. This is due to the existence of fast encoding and decoding algorithms.

Definition

A linear code C is a **cyclic code** if C satisfies:

$$(c_1, c_2, \dots, c_{n-1}, c_n) \in C \Rightarrow (c_n, c_1, \dots, c_{n-1}) \in C,$$

For every $c \in \mathbb{F}_q^n$.

Table of Contents

History

Group-Ring Codes

Code-Checkable Group Rings

References

Group Rings

Definition

Given a group G and a ring R the **group ring** RG is the ring consisting of the set of all formal finite sums $\sum_{g \in G} \alpha_g g$, where $\alpha_g \in R$.

For $u = \sum_{g \in G} \alpha_g g$, $v = \sum_{g \in G} \beta_g g \in RG$ and $\alpha \in R$, define:

- ▶ $u + v = \sum_{g \in G} (\alpha_g + \beta_g) g$,
- ▶ $uv = (\sum_{g \in G} \alpha_g g)(\sum_{h \in G} \beta_h h) = \sum_{g \in G} (\sum_{h \in G} \alpha_h \beta_{h^{-1}g}) g$,
- ▶ $\alpha u = \sum_{g \in G} (\alpha \alpha_g) g$.

Basic Properties of Group Rings

- ▶ The group ring RG is a ring.
- ▶ The group ring RG is a left R -module.
- ▶ When R is a field, the group ring RG is an algebra over R and it is called group algebra.

Theorem

For a fixed listing of elements of a finite group $G = \{g_1, g_2, \dots, g_n\}$ there is a one-to-one correspondence between RG and a subring of the matrix ring $M_n(R)$, given by:

$$w = \sum_{i=1}^n \alpha_i g_i \rightarrow W = \begin{bmatrix} \alpha_{g_1^{-1}g_1} & \alpha_{g_1^{-1}g_2} & \cdots & \alpha_{g_1^{-1}g_n} \\ \vdots & \vdots & \ddots & \vdots \\ \alpha_{g_n^{-1}g_1} & \alpha_{g_n^{-1}g_2} & \cdots & \alpha_{g_n^{-1}g_n} \end{bmatrix}$$

Group-Ring Codes

Let RG be a group ring, W a submodule of RG and $u \in RG$.

Definition

- ▶ A **right group ring encoding** is a mapping $f : W \rightarrow RG$, such that $f(x) = xu$.
- ▶ The **group-ring code** C generated by u relative to W is the image of the group ring encoding: $C = \{xu : x \in W\}$.

If u is a zero-divisor (resp. unit), C is called zero-divisor (resp. unit-derived) code.

Table of Contents

History

Group-Ring Codes

Code-Checkable Group Rings

References

Checkable Codes

- ▶ Suppose that u is a zero-divisor in RG and let v be a non-zero element such that $uv = 0$.
- ▶ Let $C = \{xu : x \in W\} = Wu$ be a code generated by u relative to W . Then

$$y \in C \Rightarrow yv = 0.$$

- ▶ If the zero-divisor code C satisfies: $y \in C \Leftrightarrow yv = 0$. Then C is called **checkable code**. In other words,

$$C = \{y \in RG : yv = 0\}.$$

Code-Checkable Group Rings

C is said to be checkable if $C = \{y \in RG : yv = 0\}$ for some $v \in RG$.

Definition (Jitman, 2010)

A group ring RG is said to be **code-checkable** if every ideal in RG is a checkable code.

Characterization of Code-Checkable Group Algebras

Let G be a finite abelian group and \mathbb{F} be a finite field of characteristic p .

Proposition (Jitman, 2010)

The group algebra $\mathbb{F}G$ is code-checkable if and only if it is a PIR.

Theorem (Fisher and Sehgal, 1976)

The group algebra $\mathbb{F}G$ is a PIR if and only if a Sylow p -subgroup of G is cyclic.

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Definition

Let π be a finite set of primes. A finite group G is called π' -by-cyclic π , if there is a normal subgroup $H \triangleleft G$ such that:

- ▶ $|H|$ is coprime with each prime in π .
- ▶ The quotient group G/H is cyclic and a π -group.

Example

Let $\pi = \{2\}$. Since $A_3 \triangleleft S_3$, $|A_3| = 3$ and $|S_3/A_3| = 2$. Then S_3 is π' -by-cyclic π .

Characterization of Code-Checkable Group Rings

Lemma

Let R be a finite commutative ring and G a finite group. Then RG is code-checkable if and only if RG is a principal ideal group ring.

Theorem (Dorsey, 2006)

Let R be a finite semisimple ring and G a finite group. Then RG is PIR if and only if G is π' -by-cyclic π , where π is the set of noninvertible primes in R .

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Characterization of Code-Checkable Group Rings

Theorem

Let G be a finite group, R a finite commutative semisimple ring and π the set of noninvertible primes in R . Then the group ring RG is code-checkable if and only if G is π' -by-cyclic π .

Table of Contents

History

Group-Ring Codes

Code-Checkable Group Rings

References

References I

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Thank You for Your Time.